

# Linear Optimization

Today we start our last topic of the semester, linear optimization.

## Comprehension goals:

- ▶ What is a linear program?
- ▶ Visualizing linear programs graphically.
- ▶ Understanding solutions graphically.
- ▶ Solving linear programs using *Mathematica*

## Fertilizer example (p.253)

A fertilizer manufacturer uses nitrates and phosphates to make batches of two different kinds of fertilizer.

- ▶ Sod-King fertilizer needs 4 phosphates, 18 nitrates.
- ▶ Gro-Turf fertilizer needs 1 phosphate, 15 nitrates.

The profit for one batch of Sod-King is \$1000.

The profit for one batch of Gro-Turf is \$500.

The company has 10 phosphates and 66 nitrates on hand.

*Question.* How many batches of each should the company make to earn the most profit?

*Initial thoughts?*

## Fertilizer example (p.253)

Translate the problem into mathematics:

We must determine how many batches to make of each.

- ▶ Let  $x$  represent the number of batches of Sod-King made.
- ▶ Let  $y$  represent the number of batches of Gro-Turf made.

*What are the constraints on what  $x$  and  $y$  can be?*

- ▶ Phosphate constraint:
- ▶ Nitrate constraint:
- ▶ Non-negativity constraints:

*What are we trying to maximize?*

- ▶ Profit:

# Linear Programs

Maximize  $1000x + 500y$

subject to  $4x + y \leq 10$

the constraints:  $18x + 15y \leq 66$

$x \geq 0$

$y \geq 0$

This is a **linear program**, an optimization problem of the form:

Maximize  $c_1x_1 + c_2x_2 + \cdots + c_nx_n$  (*the objective function*)

subject to  $a_{11}x_1 + a_{12}x_2 + \cdots + a_{1n}x_n \leq b_1$

(*the constraints*):  $a_{21}x_1 + a_{22}x_2 + \cdots + a_{2n}x_n \leq b_2$

$\vdots$

$\vdots$

$a_{m1}x_1 + a_{m2}x_2 + \cdots + a_{mn}x_n \leq b_m$

# Linear Programs

Maximize  $c_1x_1 + c_2x_2 + \cdots + c_nx_n$  (*the objective function*)

subject to

$$a_{11}x_1 + a_{12}x_2 + \cdots + a_{1n}x_n \leq b_1$$

(*the constraints*):

$$a_{21}x_1 + a_{22}x_2 + \cdots + a_{2n}x_n \leq b_2$$

$$\vdots$$

$$a_{m1}x_1 + a_{m2}x_2 + \cdots + a_{mn}x_n \leq b_m$$

Notes about linear programs:

- ▶ Constraints may be of the form  $\leq$ ,  $=$ , or  $\geq$ .
- ▶ The  $x_i$  variables are called **decision variables**.
- ▶ The decision variables can have any real value, not only integers.
- ▶ All constraints and the objective functions are *linear combinations* of the decision variables. (Coefficients are constants.)
- ▶ A linear program in the above form is “easy to solve”.

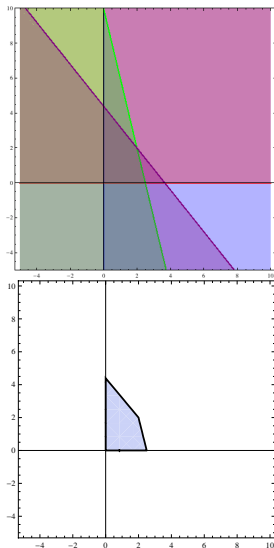
# Fertilizer example, graphically

$$\begin{array}{ll}
 \text{Maximize } 1000x + 500y & \\
 \text{subject to} & 4x + y \leq 10 \\
 \text{the constraints:} & 18x + 15y \leq 66 \\
 & x \geq 0 \\
 & y \geq 0
 \end{array}$$

Let's consider our example graphically.

*Definition:* The set of points  $(x, y)$  that satisfy the constraints is called the **feasible region**.

- ▶ In general, points of form  $(x_1, x_2, \dots, x_n)$ .
- ▶ Feasible region always a polytope. (Always has flat sides and is convex.)
- ▶ Feasible region may be bounded or unbounded; might be empty.



# Fertilizer example, graphically

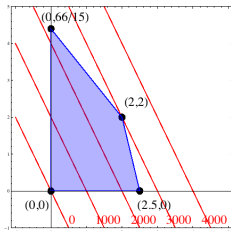
Maximize  $1000x + 500y$

subject to  $4x + y \leq 10$

the constraints:  $18x + 15y \leq 66$

$x \geq 0$

$y \geq 0$



★ The solution to the optimization problem will be the point in the feasible region that optimizes the objective function. ★

Is there a point in the feasible region such that  $1000x + 500y = 2000$ ?

Is there a point in the feasible region such that  $1000x + 500y = 4000$ ?

As we plot these **constant-objective** lines, we notice that

- ▶ They are parallel.
- ▶ If there is a feasible region, at least one line will intersect it.
- ▶ As we increase the “constant”, the last place we touch the feasible region is \_\_\_\_\_.

# Linear Optimization

We have intuited the following theorem.

*Theorem.* The maximum (or minimum) in a linear program either:

- 1 Doesn't exist (then we call the problem unbounded)
- 2 Occurs at a corner point of the feasible region.

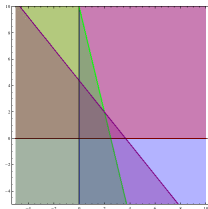
Strategy for solving a linear optimization problem:

- 0 Determine the decision variables, objective function, and constraints.
- 1 Draw the feasible region.
- 2 Compute the coordinates of all corner points.
- 3 Evaluate the objective function at each corner point.
- 4 Pick out the optimum value.



# Solution of fertilizer example

$$\begin{aligned} \text{Maximize } p(x, y) &= 1000x + 500y \\ \text{subject to} \quad & 4x + y \leq 10 \\ \text{the constraints: } & 18x + 15y \leq 66 \\ & x \geq 0 \\ & y \geq 0 \end{aligned}$$



- 1 Draw the feasible region. (Done!)
- 2 Compute the coordinates of all corner points.
  - ▶ Find the constraints that intersect; solve the associated equalities.
  - ▶  $x \geq 0$  and  $y \geq 0$ :  $(0, 0)$ . **(Not all intersections!)**
  - ▶  $x \geq 0$  and  $18x + 15y \leq 66$ :  $(0, 22/5)$ .
  - ▶  $y \geq 0$  and  $4x + y \leq 10$ :  $(5/2, 0)$ .
  - ▶  $18x + 15y \leq 66$  and  $4x + y \leq 10$ :  $(2, 2)$ .
- 3 Evaluate the objective function at each corner point.
  - ▶  $p(0, 0) = 0$                       ▶  $p(0, 22/5) = 2200$
  - ▶  $p(5/2, 0) = 2500$               ▶  $p(2, 2) = 3000$ .
- 4 Pick out the optimum value. [Max value: \$3000, occurs at  $(2, 2)$ .]

## Using *Mathematica* to solve a linear program

Once you have written your optimization problem as a linear program, you can use *Mathematica* to solve your problem.

Use either the `Maximize` or `Minimize` command.

Syntax: `Maximize[{obj, constr}, vars]`

- ▶ *obj* is the objective function that you wish to optimize.
- ▶ *constr* are the set of all constraints, joined with `&&`'s (ANDs).
- ▶ *vars* is the set of variables.

```
In[1]: Maximize[{1000 x + 500 y,  
               x >= 0 && y >= 0 && 4 x + y <= 10 && 18 x + 15 y <= 66}, {x, y}]
```

```
Out[1]: {3000, {x -> 2, y -> 2}}
```

The output gives the optimum value and the values the variables take on there.